



Improved Key Management for Fiat-Tassa Dynamic Traitor Tracing

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Introduction

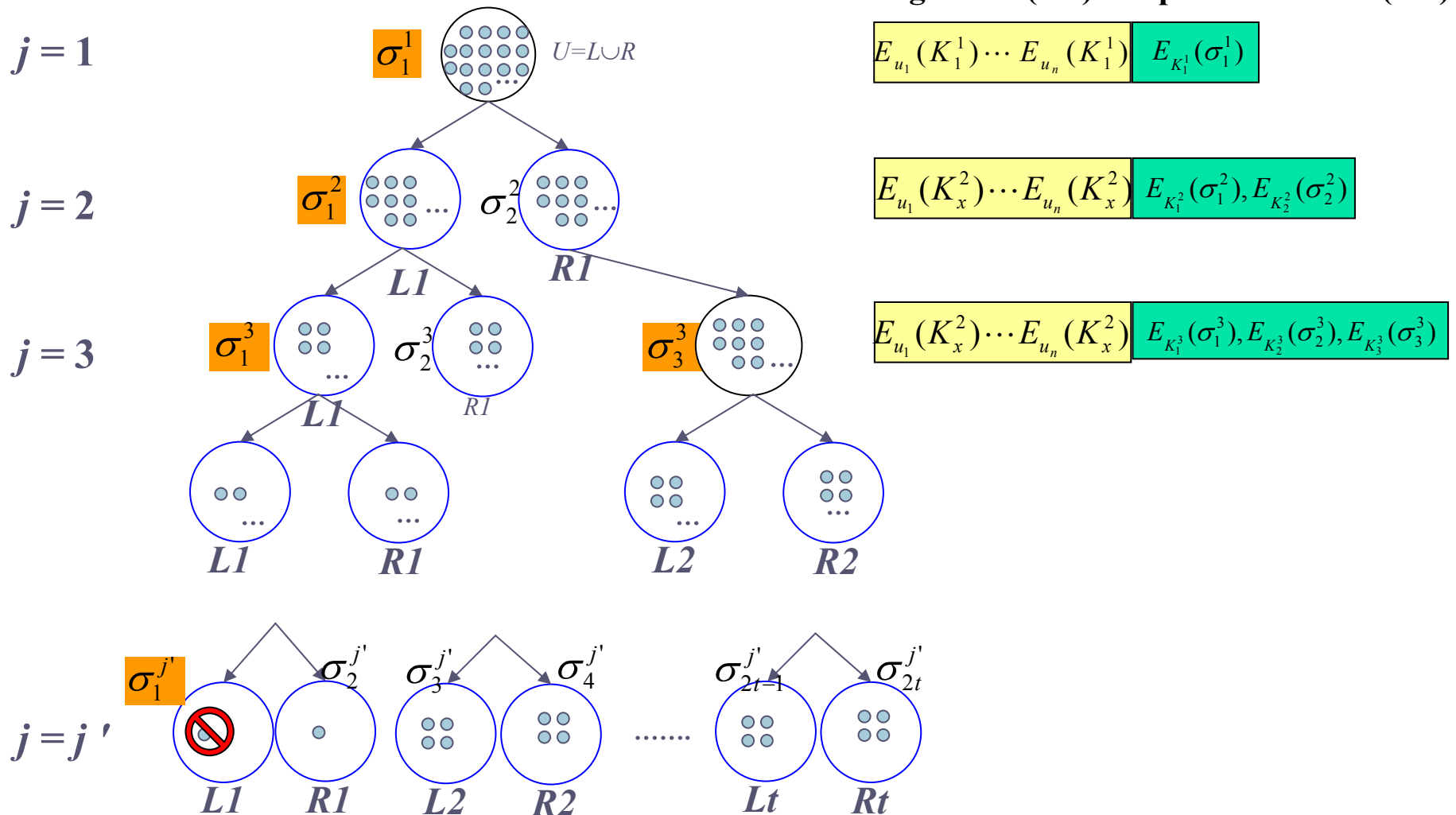
- Fiat and Tassa (CRYPTO 1999) proposed a new traitor tracing mechanism by using fingerprinting techniques that guarantee to identify at least one traitor when subscribers collude to construct a pirate decoder
- In the Fiat-Tassa scheme, each subscriber has only one personal key initially, but the broadcaster should need to construct an enabling block that contains $O(n)$ encrypted keys used for recovering the session key



Introduction (cont.)

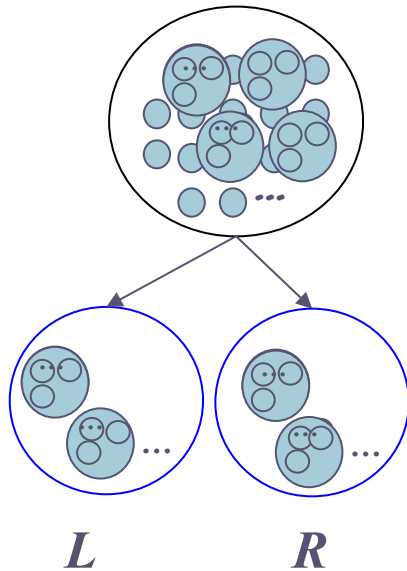
- The Fiat-Tassa scheme requires a large amount of transmission bandwidth and computational overhead for broadcast encryption
- Our proposal uses the Subset-Difference method proposed by Naor et al . (CRYPTO 2001) to construct a key tree hierarchy formed by the personal keys for the subscribers
- The improved key management mechanism preserves the security properties provided by the original Fiat-Tassa scheme
- The size of the enabling block produced by our proposal is reduced significantly from our experimental results



Tracing Process in Fiat-Tassa Scheme



Main Idea of Our Proposal

- Divide the set of n subscribers into smaller disjoint subsets
- All subscribers in the same subset share a set of secret keys



 subscriber
 subscriber subset S_{x_i, y_i}

Enabling Block (EB)

$$E_{u_1}(K_1^1) \cdots E_{u_n}(K_1^1)$$

Fiat-Tsaas

n

$$E_{S_{x,y}}(K_1^1) \cdots E_{S_{x',y'}}(K_1^1)$$

Our proposal

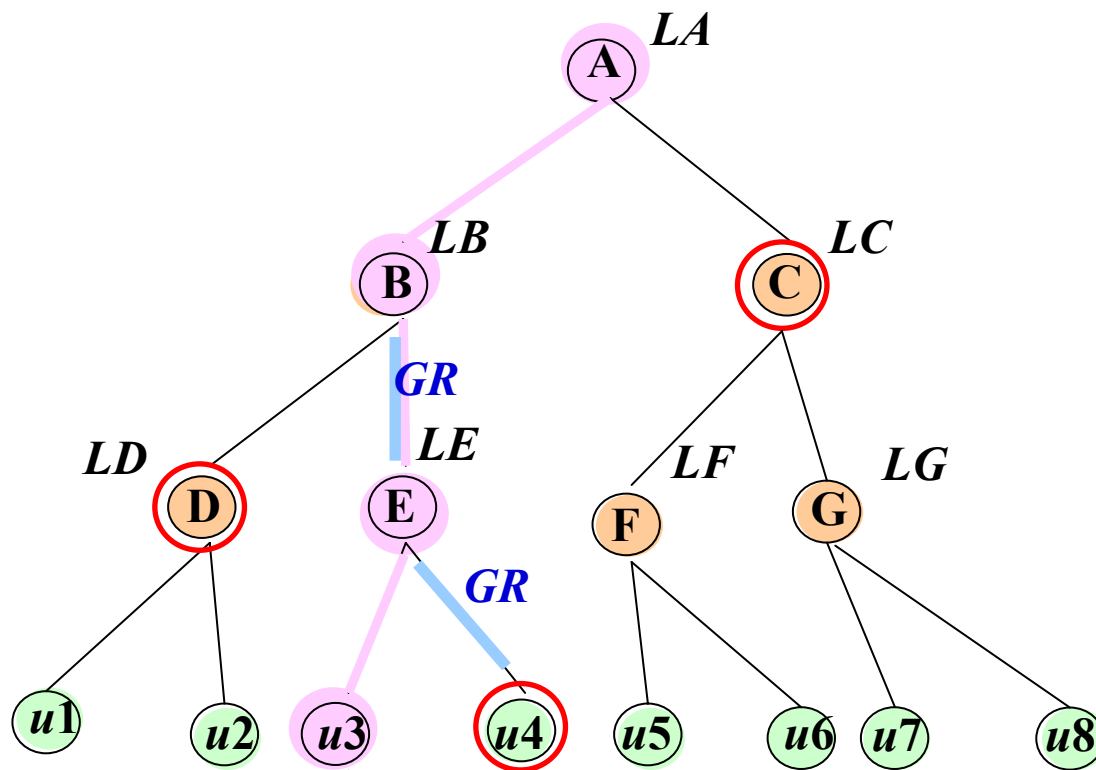
c

$$c = F(t, spt)$$

t : # of traitors could be detected

spt : pre-defined size of subset

Key Hierarchy of Our Proposal using Subset-Difference method



$L_{A,C}, L_{A,D}, L_{A,u4},$
 $L_{B,D}, L_{B,u4},$
 $L_{E,u4},$
initial

GR: key generation function

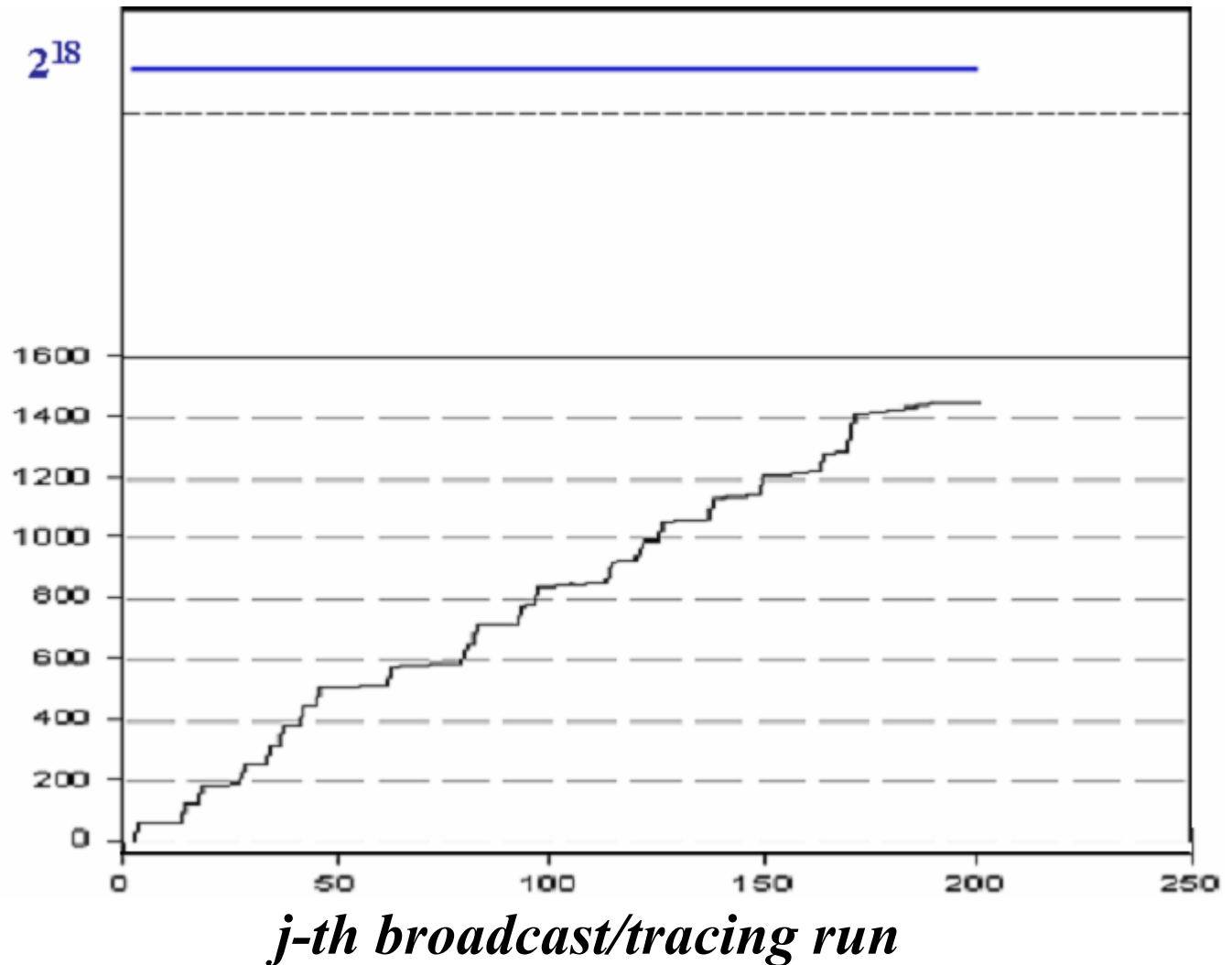
$$L_{B,u4} = GR(GR(LB))$$

$$L_{E,u4} = GR(LE)$$

Experimental Result of

Our Proposal with $n=2^{18}$, $spt=64$

EB size c



Fiat-Tassa scheme

Our proposal



The First Glimpse of Comparison

	Fiat-Tassa	Our Proposal
Personal keys	1	$\log^2 n$
Enabling block size	n	$c = F(t, spt)$
Rounds for detecting traitor	$\log(n)$	$\log(n/spt)$

Notes:

t : # traitors could be detected in each round

spt : pre-defined size of subset



Thank you

for your attention